

# Lecture 07

## Earlier...

**1 Theorem.** If  $\text{Disc}_\mu(F) \leq 2^{-V}$ , for  $v \gg (\log n)^2$ , then any majority of  $s$  linear threshold functions to compute  $F$  must have  $s = \frac{2^{\Omega(V)}}{n}$ .

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**2 Gordan's Transposition Theorem.** Let  $M \in \mathbb{R}^{m \times n}$ . Then exactly one of the following statements must hold:

- (i) For some vector  $u$ ,  $Mu \geq \varepsilon$  (all coordinates of  $Mu$  are positive);
- (ii)  $M^T v = 0$  for some non-zero vector  $v \geq 0$ .

*Proof.* Let  $e, e'$  be all-1 vectors of suitable dimension. Consider the following dual pair of linear programs:

$$P \left\{ \begin{array}{ll} \min & \langle e, z \rangle \\ \text{s.t.} & Mu + z \geq e' \\ & z \geq 0 \end{array} \right. \quad Q \left\{ \begin{array}{ll} \max & \langle e', v \rangle \\ \text{s.t.} & M^T v = 0 \\ & v \leq e \\ & v \geq 0 \end{array} \right.$$

Clearly both  $P$  and  $Q$  are feasible;  $z = e'$  satisfies  $P$ , and  $v = 0$  satisfies  $Q$ . Hence both programs have solutions, regardless of  $M$ .

Now suppose that for any  $u$ , some coordinate of  $Mu$  is  $\leq 0$ . Then any  $z$  satisfying the constraints of  $P$  must have some coordinate  $> 0$ . This is also true for the optimum solution  $u^*, z^*$  of  $P$ , and hence the optimum value of  $P$  is  $> 0$ . But this equals the optimum value of  $Q$ , and so the optimum solution  $v = v^*$  of  $Q$  must be non-zero (and have  $M^T v = 0$ ).

Now suppose instead that some  $u$  has  $Mu \geq \varepsilon$ ; then we may assume  $\varepsilon = 2$  by multiplying  $u$  with a sufficiently large number. Then  $z = 0$  is a solution to  $P$ , and so 0 is the optimum value of both  $P$  and  $Q$ . But that implies that  $v = 0$  is the only solution to  $Q$ . ■

**3 Corollary.** For any  $f, d$ , exactly one of the following holds:

(i)  $\deg_{\pm}(f) \leq d$ ;

(ii) there exists  $\rho$  such that  $\sum_x \rho(x) f(x) \chi_S(x) = 0$  whenever  $|S| \leq d$ .

(proof of thm)

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## Degree-discrepancy Theorem

The following upper bound on discrepancy will be used

**4 Lemma (Discrepancy Bound).**  $f : X \times Y \rightarrow \{1, -1\}$ ,  $\mu \sim X \times Y$ .

$$\text{Disc}_{\mu}(f)^2 \leq |X| \sum_{y, y' \in Y} \left| \sum_{x \in X} \mu(x, y) \mu(x, y') f(x, y) f(x, y') \right|$$

The proof in the original paper is sufficiently simple and straightforward. We will not copy it here.

**5 Theorem (Degree/Discrepancy theorem).** Let  $\deg_{\pm}(f) = d \geq 1$ . Let  $N \geq n$ , and define  $F(x, S) = f(x|_S)$  with  $x \in \{1, -1\}^N$  and  $S \in \binom{[N]}{n}$ . Then

$$\text{Disc}_{\mu}(F) \leq \left( \frac{4en^2}{Nd} \right)^{d/2} \quad (\leq 2^{-d/2} \text{ for } N \geq n^2/d).$$

*Proof of Degree/Discrepancy theorem.* For any given distribution  $\lambda$  on  $x, S$ , the Discrepancy bound lemma says that:

$$\text{Disc}_\lambda(f)^2 \leq 2^N \sum_{S,T} \left| \sum_x \lambda(x, S) \lambda(x, T) f(x|_S) f(x|_T) \right|.$$

So let  $\mu$  be given by §3, and

$$\lambda(x, S) = 2^{-N+n} \binom{N}{n}^{-1} \mu(x|_S),$$

i.e.,  $\lambda$  first picks  $S$  uniformly, then chooses  $x|_S$  according to  $\mu$ , and picks the remaining  $x$ -coordinates uniformly.

Then,

$$\begin{aligned} \text{Disc}_\lambda(f)^2 &\leq 2^N \sum_{S,T} \binom{N}{n}^{-2} \left| \sum_x 2^{-2N+2n} \mu(x, S) \mu(x, T) f(x|_S) f(x|_T) \right| \\ &= 2^{2n} E_{S,T \sim \mathcal{U}} [|\Gamma(S, T)|], \end{aligned}$$

where  $\Gamma(S, T) = E_{x \sim \mathcal{U}} [\mu(x, S) \mu(x, T) f(x|_S) f(x|_T)]$ .

$\Gamma(S, T)$  is sort-of a measure of  $f$ -correlation among  $S$  and  $T$  wrt  $\mu$ . If  $|S \cap T| = k$ , then regardless of the distribution  $\mu$ , we can at least say that  $|\Gamma(S, T)| \leq 2^{k-2n}$ . Indeed, suppose that  $S = \{x_1, \dots, x_n\}$  and  $T = \{x_1, \dots, x_k, x_{n+1}, \dots, x_{2n-k}\}$ . Then:

$$\begin{aligned} |\Gamma(S, T)| &\leq E_{x \sim \mathcal{U}} |\mu(x|_S) \mu(x|_T) f(x_S) f(x_T)| \\ &= \sum_x 2^{-N} \mu(x|_S) \mu(x|_T) \\ &= \sum_{x|_S} \left( 2^{-n} \mu(x|_S) \sum_{x|_{T \setminus S}} 2^{-n+k} \mu(x|_T) \right) \\ &\leq \left( \sum_{x|_S} 2^{-n} \mu(x|_S) \right) \left( \max_{x_1, \dots, x_k} \sum_{x|_{T \setminus S}} 2^{-n+k} \mu(x|_T) \right) \\ &\leq 2^{-n} \cdot 2^{-n+k} \end{aligned}$$

(first inequality is the triangle inequality, and the last follows because  $\mu$  is a distribution).

But when  $|S \cap T| = k < d$ , we can actually show that  $\Gamma(S, T) = 0$ , for the special distribution  $\mu$  of §3. Indeed, assume  $S$  and  $T$  are given as before; then

$$\begin{aligned} \Gamma(S, T) &= \sum_x 2^{-N} \mu(x_1, \dots, x_n) f(x_1, \dots, x_n) \mu(x|_T) f(x|_T) \\ &= \sum_{x_1, \dots, x_n} \mu(x_1, \dots, x_n) f(x_1, \dots, x_n) \left( \sum_{x|_{[N] \setminus [n]}} 2^{-N} \mu(x|_T) f(x|_T) \right) \end{aligned}$$

The right-hand parenthesis is actually a function  $g$  which depends only on  $< d$  (boolean) coordinates  $\{x_1, \dots, x_k\}$ , and the whole expression is the measure of  $f$ -correlation between  $f$  and  $g$  with respect to  $\mu$ ; which equals zero, by our choice of  $\mu$ .